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Ce cahier a également été publié par le Centre interuniversitaire de recherche en économie quantitative (CIREQ) sous le numéro 07-2008.

This working paper was also published by the Center for Interuniversity Research in Quantitative Economics (CIREQ), under number 07-2008.

ISSN 0709-9231

Priorities in the Location of Multiple Public Facilities*

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February 2008

Abstract

A collective decision problem is described by a set of agents, a profile of single-peaked preferences over the real line and a number k of public facilities to be located. We consider public facilities that do not suffer from congestion and are non-excludable. We provide a characterization of the class of rules satisfying *Pareto-efficiency*, *object-population monotonicity* and *sovereignty*. Each rule in the class is a *priority rule* that selects locations according to a predetermined priority ordering among “interest groups”. We characterize each of the subclasses of priority rules that respectively satisfy *anonymity*, *hiding-proofness* and *strategy-proofness*. In particular, we prove that a priority rule is strategy-proof if and only if it partitions the set of agents into a fixed hierarchy. Alternatively, any such rule can be viewed as a collection of fixed-populations *generalized peak-selection median rules* (Moulin, 1980), that are linked across populations, in a way that we describe.

Keywords: Multiple public facilities, Priority rules, Hierarchical rules, Object-population-monotonicity, Sovereignty, Anonymity, Strategy-proofness, Generalized median rules, Hiding-proofness.

JEL classification codes: D60; D63; D70; D71; H41.

*We thank Anna Bogomolnaia, Haluk Ergin, Lars Ehlers, Bettina Klaus, François Maniquet, Hervé Moulin, Jorge L. Garcia Ramirez, Yves Sprumont and William Thomson for helpful conversations or comments. Part of this paper was written when the first author visited the Université de Montréal, and when the second author visited Maastricht University. Financial support from CIREQ and METEOR is gratefully acknowledged.

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1 Introduction

We consider a generalization of the unidimensional voting model studied by Black (1948) and Moulin (1980). A collective decision problem is described by a set of agents, a profile of single-peaked preferences over the real line, and a number k of public facilities to be located –e.g. public libraries. Each public facility is non-excludable and does not exhibit congestion. Thus, it is a pure public good. Each agent has preferences over sets of locations that reflect the fact that she will use only her preferred facility among the k that are provided. We search for a rule that decides on how to locate the facilities, in any possible collective decision problem. We follow the axiomatic method. In addition to *Pareto-efficiency*, we search for rules that satisfy two normative properties, which we introduce in this paper, *object-population monotonicity* and *sovereignty*.

“Public facilities” should be interpreted in a broad sense. They could be actual facilities, such as wireless towers, to be located in space. They could be varieties of software to be provided to a community of users. In this example, the “locations” represent different program varieties. They could also be the official languages for a linguistically diversified society. We will use this third example, to explain the content of our axioms in a concrete manner.

The European Union (henceforth, EU) has to select a set of official working languages for its main tribune, among the set of all the languages used in EU countries.¹ Not all these languages can be selected as official working language. How should the EU decide which languages to select as official working ones? In addition, how should it modify this selection after an enlargement? We would like to have a method for selecting official working languages that respects a

¹Currently, the official working languages of the EU are English, French and German. For a more detailed discussion concerning the choice of official languages in the EU, see the paper by Fidrmuc, Ginsburgh and Weber (2006).

few desirable properties.² As a first requirement, we require *Pareto-efficiency*. The selected set of languages should be such that no other set with the same cardinality is weakly preferred by all EU citizens and strictly preferred by at least one citizen. In addition, we are interested in two new properties –*object-population monotonicity* and *sovereignty*. Roughly speaking, *object-population monotonicity* says that if a new group of citizens joins the EU, and at the same time, the number of official working languages increases, in such a way that the increase of the number of languages compensates the increase in linguistic diversity in the EU due to the arrival of new citizens, then all agents initially present should weakly benefit from the change. Last, suppose that a single language needs to be selected as communicating language within EU institutions. *Sovereignty* loosely says that any language could be selected as the unique working language, provided that an appropriately selected, and possibly large, “group of interest” defending this particular language, is brought into the EU.

Equipped with these axioms, we provide a complete characterization of the class of rules that jointly satisfy them. A *priority rule* operates as follows. First, it partitions the EU population into linguistic groups, i.e. groups of citizens that share the same preferred language. Second, it ranks these linguistic groups according to some predetermined priority order that can depend on the identities of their members and their full preferences over all languages. The rule then selects the preferred languages of the top k groups in the priority order. Our main result is that any rule that satisfies our three axioms is a priority rule (Theorem 1). All priority rules are *efficient*, and *object-population monotonic*, but not all are *sovereign*. We characterize the class of priority rules that satisfy all three properties (Theorem 1). The corresponding set is large within the set of priority rules, in the following sense: Any priority rule can be approximated

²In the process of its enlargement, the EU may repeatedly face the choice of adding new languages. Hence, we need a method –a rule in the language of collective choice– that can be used in different situations.

with a sequence of rules that satisfy all three properties (Theorem 2). Priority rules form a rich class. In particular, it includes an interesting subclass of *majoritarian rules* that rank linguistic groups according to their cardinality.

Next, we investigate on priority rules that satisfy additional desirable properties, such as *anonymity*, *hiding-proofness*, *median-selection* and *strategy-proofness*. Unfortunately, no priority rule selects the median voter when a single language needs to be selected and the number of agents is odd. The priority rules that are *strategy-proof* form an interesting subclass of *hierarchical rules* (Theorem 3). Each such rule partitions the EU population into a fixed hierarchy of priority levels. Each such rule can also be described as a linked collection of generalized peak-selecting median rules (Moulin, 1980) restricted in a particular way. We provide a complete description of the linkage across populations imposed by our axioms on such a collection.

The problem of locating a single facility is well-studied in the voting literature. Moulin (1980), Ching (1997), Barberà and Jackson (1994) among others, have studied the strategic properties of rules for locating a single facility. Other scholars have studied rules for this problem, that satisfy normative properties. The principle of solidarity says that when circumstances change, all agents not responsible for the change should be affected in the same direction.³ Thomson (1993), Ching and Thomson (1996), Vohra (1998) and Klaus (2001) investigate the two main formulations in this context.⁴ In all the papers cited in this para-

³Thomson (1995, 1999a, 1999b) provides excellent surveys on the different formulations of the solidarity principle and its applications to different economic contexts. *Replacement-dominance* operates on preference profiles for a fixed population. It says that when the preferences of one agent change, all the other agents, whose preferences are kept fixed, should be affected in the same direction. *Population-monotonicity* operates on preference profiles for a variable population. It says that when new agents join the economy, all the agents who were initially present should all be affected in the same direction.

⁴Gordon (2007b) studies the solidarity principle in the location of a public facility on a cycle. Gordon (2007b) generalizes this literature on the solidarity principle to a more abstract non-geometric setting that contains location problems of single and multiple facilities as special cases.

graph, the rules select the location of a single facility, hence resources are kept fixed.

The problem of locating two facilities was studied first by Miyagawa (1998, 2001). In these works, Miyagawa proposes to extend preferences over single locations to finite sets of locations according to the *max-extension*. This extension compares finite sets of locations by comparing only their best element. Ehlers (2002, 2003) also studies this problem, but extends preferences in a *lexicographic* manner. Both Miyagawa (1998, 2001) and Ehlers (2002, 2003) follow the normative route. They look for rules that satisfy either *replacement-domination* or *population-monotonicity*. But once again, the set of resources is kept fixed, since their rules always locate exactly two facilities. In contrast, we do not restrict the set of facilities to any particular cardinality. Thus, our rules prescribe how to locate any number of facilities. In addition, our *object-population monotonicity* axiom operates across different populations and across problems with different numbers of facilities. It can be viewed as a formulation of the solidarity principle in a context where both population and resources can vary. Our work resembles the work of these scholars in style, in the sense that we also characterize families of rules using axioms, but the axioms we study are different, and the set of location problems we cover is larger.

Barberà and Beviá (2002, 2005) and Ju (2007) also study the location of multiple facilities. They focus on *consistency*. This property says that when one $k - 1$ facility and its “users” are simultaneously removed, the remaining location should remain unchanged. These authors show that the set of rules that satisfy this property and additional interesting requirements is large, but they do not provide a characterization. In contrast, we obtain several characterizations on the basis of *object-population monotonicity* and other axioms.

Our main contribution is that we identify and axiomatize the class of priority rules, in an environment for which few rules have been proposed. These rules

have a simple parametric structure. The structure, in turn, facilitates further analysis: It is relatively easy to characterize the set of priority rules that satisfy an additional axiom. To obtain our main characterization, we introduce two new axioms. These axioms are meaningful in the context of the location of public facilities, but they also are of independent interest, and could be further investigated in other contexts.

The plan of the paper is as follows. In Section 2, we introduce the model and necessary definitions. In Section 3, we introduce priorities and priority rules. In Section 4, we present various examples of priorities and priority rules. In Section 5, we present our main axioms. In Section 6, we present our main characterization. In Section 7, we show that the set of rules we characterize is dense within the set of the priority rules for an appropriately chosen topology. In Section 8, we study the subclasses of priority rules that satisfy respectively anonymity, and hiding-proofness. We also show that no priority rule selects the median in problems with one good and a population with an odd cardinality. In Section 9, we study the subclass of strategy-proof priority rules.

2 The model

The following definitions are useful throughout the paper. A binary relation over a set Z is a subset of Z^2 . A binary relation r over Z is transitive if for all $x, y, z \in Z$, we have $(x r y \text{ and } y r z) \implies (x r z)$. It is reflexive if for all $x \in Z$, we have $x r x$. It is complete if for all $x, y \in Z$, we have either $x r y$ or $y r x$. It is symmetric if for all $x, y \in Z$, we have $(x r y) \implies (y r x)$. It is asymmetric if for all $x, y \in Z$, we have $(x r y) \implies \neg(y r x)$. A binary relation over Z is a *weak ordering* over Z if it is transitive, reflexive and complete. It is a *strict ordering* over Z if it is transitive, asymmetric and complete. It is a *weak partial order* if it is transitive and reflexive. It is an *indifference relation* if it is transitive, reflexive, and symmetric. We will later introduce other types

of binary relations.

There is a countably infinite set \mathbb{N} of potential agents. A *population* N is a finite and nonempty subset of \mathbb{N} . The population is collectively endowed with a number k of identical public facilities, each to be located on the real line \mathbb{R} . A typical location on \mathbb{R} is denoted by x . An *assignment* is a menu of locations, i.e. a finite subset $X \subset \mathbb{R}$. A *k-assignment* is an assignment for exactly k facilities, i.e. a subset $X \subset \mathbb{R}$ such that $|X| = k$. Let \mathcal{X}_k be the class of all k -assignments. In particular, a 1-assignment is a single location $x \in \mathbb{R}$, so that $\mathcal{X}_1 = \mathbb{R}$. Let $\mathcal{X} \equiv \cup_{k \geq 1} \mathcal{X}_k$ be the class of all assignments.

A *preference* over \mathcal{X} is a weak ordering over \mathcal{X} . Each agent $i \in N$ has a preference R_i over \mathcal{X} . For each preference R_i , let P_i and I_i stand for the strict ordering and the indifference relation associated with R_i , respectively. We restrict attention to the class \mathcal{R} of *single-peaked* preferences over \mathcal{X} , defined by the following two conditions. The first condition is the common single-peakedness notion, for preferences over single locations on the real line. The second condition extends preferences from single locations to menus.⁵ A preference R_i is single-peaked if the following holds.

- i) There is a location $p(R_i)$, such that for all $x, y \in \mathbb{R}$ satisfying either $x < y \leq p(R_i)$ or $p(R_i) \geq y > x$, we have $y P_i x$. The location $p(R_i)$ is called the *peak* of preference R_i .
- ii) For all $X, Y \in \mathcal{X}$, we let $X R_i Y$ if there is $x \in X$ such that for all $y \in Y$, we have $x R_i y$.

For each population N , a *preference profile for* N is a list $R_N = (R_i)_{i \in N} \in \mathcal{R}^N$. More generally, a *preference profile* is a preference profile for some popu-

⁵There are different ways to extend preferences over points to preferences over sets. Consistent with the definition of a public facility used in this paper, we consider the max-extension of preferences used by Miyagawa (2001).

lation N .⁶ For each profile R_N and each subpopulation $K \subseteq N$, let R_K denote the subprofile $(R_i)_{i \in K}$. For each profile $R_N \in \mathcal{R}^N$, let $p(R_N)$ be the set of *peak locations* for R_N , i.e. $p(R_N) \equiv \{p(R_i) : i \in N\}$. For each $k > 0$, let \mathcal{P}_k be the set of preference profiles R_N with a number of distinct peak locations greater than or equal to k , i.e. such that $k \leq |p(R_N)|$. A problem is a pair (k, R_N) such that k is a positive integer, and $R_N \in \mathcal{P}_k$.⁷

A *rule* is a sequence $f = \{f_1, f_2, \dots\}$ of mappings $f_k : \mathcal{P}_k \rightarrow \mathcal{X}_k$. For each problem (k, R_N) , the rule f prescribes an assignment in \mathcal{X}_k .⁸ For each $k \geq 1$, the set of mappings f_k is $\mathcal{X}_k^{\mathcal{P}_k}$. Therefore, the set of all rules is the set $\prod_{k=1}^{\infty} \mathcal{X}_k^{\mathcal{P}_k}$.

3 Priority rules

Let us introduce a class \mathcal{F} of rules which will play an important role in our results. A profile R_M is *peak-unanimous* if all the preferences of this profile have the same peak, i.e. $p(R_M)$ is a singleton. Let \mathcal{T} be the set of peak-unanimous profiles. For any two peak-unanimous profiles R_L and R'_M , we say that R_L and R'_M are *compatible* if they have distinct peaks and disjoint populations, i.e. $p(R_L) \neq p(R'_M)$ and $L \cap M = \emptyset$.

We now introduce the class of priorities over some nonempty subset \mathcal{S} of \mathcal{T} . Let \succ be a binary relation over \mathcal{S} . The binary relation \succ is *almost complete* if for all $R_L, R_M \in \mathcal{S}$, we have $(R_L \succ R_M \text{ or } R_M \succ R_L) \Leftrightarrow (R_L \text{ and } R_M \text{ are compatible})$.⁹ It is *almost transitive* if for all $R_K, R_L, R_M \in \mathcal{S}$, such that R_K and R_M are compatible, we have $(R_K \succ R_L \text{ and } R_L \succ R_M) \implies (R_K \succ R_M)$. The binary

⁶A profile R_N specifies a population N and each of its agents' preferences.

⁷The restriction $k \leq |p(R_N)|$ allows us to focus on non-trivial cases. When $k > |p(R_N)|$, it is possible to locate one facility at each peak location, so that the welfare of each agent is maximized. Locating the remaining facilities does not affect any agent's welfare.

⁸Our definitions rule out locating more than one facility at the same point. Under single-peaked preferences, and for the class of problems we consider, Pareto-efficiency would exclude duplication anyway.

⁹In particular, an almost complete binary relation \succ over \mathcal{T} is never reflexive.

relation \succ is a *priority* over \mathcal{S} if it is asymmetric, almost transitive and almost complete.¹⁰ For each nonempty $\mathcal{S} \subseteq \mathcal{T}$, let $\mathbb{P}_{\mathcal{S}}$ be the set of priorities over \mathcal{S} .

For each profile R_N , the peak-unanimous subprofile R_M of R_N is *maximal* if $p(R_M) \cap p(R_{N \setminus M}) = \emptyset$. Any two distinct maximal peak-unanimous subprofiles are compatible. It follows that the collection of maximal peak-unanimous subprofiles of some profile is strictly ordered by any priority (see footnote 10). We are now ready to define the family of priority rules, parametrized by the set $\mathbb{P}_{\mathcal{T}}$. For each $\succ \in \mathbb{P}_{\mathcal{T}}$, the *priority rule* f associated with \succ is defined as follows. Let (k, R_N) be an arbitrary problem. Then the priority \succ strictly ranks the maximal peak-unanimous subprofiles in the decomposition of R_N and $f_k(R_N)$ selects the peak locations of the top k maximal peak-unanimous subprofiles for \succ . In more precise terms, $f_k(R_N)$ is the k -assignment such that $f_k(R_N) \subseteq p(R_N)$, and for all two maximal peak-unanimous subprofiles R_M and R_L in R_N , if $p(R_M) \subseteq f_k(R_N)$ and $p(R_L) \not\subseteq f_k(R_N)$, then $R_M \succ R_L$. Let \mathcal{F} be the set of priority rules.

4 Examples

In this section, we provide a few examples that illustrate how large and diverse the set of priorities is. We first present some priorities that can be described in a simple way (Examples 1, 2 and 3), and then move on to more complex priorities, for which it is convenient to use a lexicographic formulation (Examples 4 to 8).

Example 1: The *left-peaks priority* \succ^{LP} is such that, for all compatible $R_M, R_N \in \mathcal{T}$, we have $R_M \succ^{LP} R_N$ if and only if $p(R_M) < p(R_N)$. The *left-peaks*

¹⁰A priority \succ is not a partial order, as it is not fully transitive. However, priorities have the following important property. The restriction of a priority \succ on any set \mathcal{S} of pairwise compatible unanimous profiles is a strict ordering. If this set is finite, the priority \succ has a greatest (or top) element in \mathcal{S} . A top element for \succ typically does not generally exist on a set of unanimous profiles whose elements are not pairwise compatible, even if it is a finite set.

rule f^{LP} is the priority rule defined by \succ^{LP} . An alternative definition of the rule f^{LP} is that it selects, for each problem (k, R_N) , the k -assignment X such that for all $i \in N$, if $p(R_i) \notin X$, then $p(R_i) > \max(X)$.

Example 2: The *right-peaks priority* \succ^{RP} is such that, for all compatible $R_M, R_N \in \mathcal{T}$, we have $R_M \succ^{RP} R_N$ if and only if $p(R_M) > p(R_N)$. The *right-peaks rule* f^{RP} is the priority rule defined by \succ^{RP} . An alternative definition of the rule f^{RP} is that it selects, for each problem (k, R_N) , the k -assignment X such that for all $i \in N$, if $p(R_i) \notin X$, then $p(R_i) < \min(X)$.¹¹

Example 3: A priority \succ^{SD} is a *serial dictatorship* if there exists a strict ordering \triangleright of all agents in \mathbb{N} such that, for all compatible $R_M, R_N \in \mathcal{T}$, we have $R_M \succ^{SD} R_N$ if and only if there exists $i \in M$, such that for all $j \in N$, we have $i \triangleright j$. For example, the strict ordering could be such that for all $i, j \in \mathbb{N}$, we have $i \triangleright j$ if and only if $i < j$.¹²

Lexicographic formulations: Some priorities are easier to describe using a lexicographic formulation. This means that the priority is determined first by a primary criterion that is not always decisive. When the primary criterion is indecisive, a secondary criterion is used, which can vary depending on the pair that is being compared.

To formalize this idea, let us introduce the following definitions. Let \succeq be a weak ordering over \mathcal{T} . Let \mathbb{T} be the set of indifference classes for \succeq in \mathcal{T} .¹³

¹¹Miyagawa (2001) showed that when $|N| \geq 4$ and $k = 2$, the only mappings $\mathcal{R}^N \rightarrow \mathcal{X}_2$ satisfying *efficiency* and *replacement-domination* are the left-peaks rule and the right-peaks rule. These solutions are also anonymous and group-strategy-proof. The left-peaks and right-peaks priorities are also trivial hierarchical rules (see Example 4) where all agents belong to the same unique indifference class.

¹²The serial dictatorships are extreme cases of hierarchical rules (see Example 4) where each indifference class is a singleton.

¹³Requiring \succeq to be a weak ordering rather than an almost complete and almost transitive (and not asymmetric) binary relation on \mathcal{T} is important here, as it allows us to refer to the

For each $T \in \mathbb{T}$, let \succ_T be a priority over T . Finally, let \succ be a priority over \mathcal{T} . Then, the priority \succ has the lexicographic formulation $(\succeq, (\succ_T)_{T \in \mathbb{T}})$ if for all compatible $R_M, R_N \in \mathcal{T}$, we have $R_M \succ R_N$ if and only if $R_M \succeq R_N$ and either (i) $\neg(R_N \succeq R_M)$ or (ii) There is an indifference class T for \succeq such that $(R_M, R_N \in T$ and $R_M \succ_T R_N)$.

It should be clear that all priorities admit many lexicographic formulations. However, there are priorities for which a particular lexicographic formulation is the most natural and convenient definition. We now present a few examples of such priorities.

Example 4: A priority \succ^H is *hierarchical* if it admits a lexicographic formulation $(\succeq, (\succ_T))$ that satisfies the following conditions. (i) There is a weak ordering \triangleright of all agents in \mathbb{N} , such that, for all $R_M, R_N \in \mathcal{T}$, we have $R_M \succeq R_N$ if and only if there exists $i \in M$, such that for all $j \in N$, we have $i \triangleright j$. (ii) For each $T \in \mathbb{T}$, the priority \succ_T is either the (restriction to T of) the left-peaks or (restriction to T of) the right-peaks priority.¹⁴

Example 5: A priority \succ^M is *majoritarian* if it admits a lexicographic formulation $(\succeq, (\succ_n))$ such that, for all $R_M, R_N \in \mathcal{T}$, we have $R_M \succeq R_N$ if and only if $|M| \geq |N|$. The indifference classes are the sets $T_n = \{R_N \in \mathcal{T} : |N| = n\}$. For each $n \geq 1$, the tie-breaking rule \succ_n can be any priority. For example, we could require \succ_n to be the left-peaks priority, for all T (*left-majoritarian priority*) or the right-peaks priority (*right-majoritarian priority*).¹⁵

indifference classes of \succeq .

¹⁴In Section 9, we show that hierarchical rules are the only priority rules that are strategy-proof. One could also define a more general class of (non strategy-proof) hierarchical rules by allowing the \succ_T to be any priority, non necessarily the left-peaks or right-peaks priorities.

¹⁵We can similarly define the family of *minoritarian* priorities by requiring instead that $R_M \succeq R_N$ iff $|M| \leq |N|$.

Example 6: A priority \succ^M is *x_0 -centralist-majoritarian* if it admits a lexicographic formulation $(\succeq, (\succ_v))$, such that there are a location $x_0 \in \mathbb{R}$ (the “center”), a distance $D(x, y)$ on \mathbb{R} and an index $u : \{1, 2, \dots\} \times \mathbb{R}_+ \rightarrow \mathbb{R}$, where $u(n, d)$ is increasing in n and decreasing in d , such that, for all $R_M, R_N \in \mathcal{T}$, we have $R_M \succeq R_N$ if and only if

$$u(|M|, D(p(R_M), x_0)) \geq u(|N|, D(p(R_N), x_0)).$$

The indifference classes are the sets

$$T_v = \{R_N \in \mathcal{T} : u(|M|, D(p(R_M), x_0)) = v\}.$$

For each such class, the tie-breaking rule \succ_v can be any priority. For example, we could require \succ_v to be the left-peaks priority for all $v \in \mathbb{R}$ (*left- x_0 -centralist-majoritarian priority*) or the right-peaks priority for all $v \in \mathbb{R}$ (*right- x_0 -centralist-majoritarian priority*).¹⁶

We conclude this section by an illustration of how the *left-majoritarian rule* f^{LM} allocates goods for two particular problems. Let $N = \{1, \dots, 10\}$, $k = 2$ and the two profiles R_N and R'_N with peak locations distributed as follows.

Profile R_N	Agents	9,10	6,7,8	1,...,5
	Peak locations	0	1	2

Profile R'_N	Agents	1,2,3	4,5,6	7,8,9	10
	Peak locations	0	1	2	3

Let R_{L_0}, R_{L_1} and R_{L_2} be the maximal peak-unanimous subprofiles of R_N , with $L_0 = \{9, 10\}$, $L_1 = \{6, 7, 8\}$, and $L_2 = \{1, \dots, 5\}$, we have $|L_0| < |L_1| < |L_2|$, therefore $R_{L_2} \succ^{LM} R_{L_1} \succ^{LM} R_{L_0}$, i.e. $f_2(R_N) = \{1, 2\}$. Similarly, let $R'_{M_0}, R'_{M_1}, R'_{M_2}$ and R'_{M_3} be the maximal peak-unanimous subprofiles of R'_N , with $M_0 = \{1, 2, 3\}$, $M_1 = \{4, 5, 6\}$, $M_2 = \{7, 8, 9\}$, and $M_3 = \{10\}$, we have $|M_0| = |M_1| = |M_2| > |M_3|$, and $p(R'_{M_0}) < p(R'_{M_1}) < p(R'_{M_2})$, thus $R'_{M_0} \succ^{LM} R'_{M_1} \succ^{LM} R'_{M_2} \succ^{LM} R'_{M_3}$, i.e. $f_2^{LM}(R'_N) = \{0, 1\}$.

¹⁶A more general family is obtained by relaxing the monotonicity assumptions on u .

5 Axioms

Our first axiom is the usual (*Pareto*)-*efficiency* axiom. For each profile $R_N \in \mathcal{R}^N$ and each $x, y \in \mathbb{R}$, we say that x weakly Pareto-dominates y for profile R_N if $x R_i y$ for each $i \in N$. This is denoted by $x R_N y$.

A rule f satisfies **efficiency** if, for each problem (k, R_N) , there is no k -assignment X such that $X R_N f_k(R_N)$, and $X P_j f_k(R_N)$ for some $j \in N$.

Any rule that always selects distinct peaks locations is efficient. Therefore, all priority rules are efficient. Next, we would like our rules to satisfy a fairness condition when resources and population vary. First, when the population is held fixed but resources vary, it is natural to require that the agents who are initially present in the economy, and whose preferences are kept fixed, all weakly gain as the number of available facilities increases.

A rule satisfies **object-monotonicity** if, for each problem (k, R_N) such that $k < |p(R_N)|$, we have $f_{k+1}(R_N) R_N f_k(R_N)$.

All priority rules are object-monotonic, since they satisfy, for each such problem, the stronger condition $f_k(R_N) \subset f_{k+1}(R_N)$. Object-monotonicity turns out to be very weak in this model. Many rules that have very little in common satisfy this axiom. This motivates looking at changes in both resources and population.¹⁷ We require that if both resources and population increase, and the new resources are sufficient to fully satisfy the newcomers, then the agents who are initially present in the economy, and whose preferences are kept fixed,

¹⁷One could also study changes in population, while resources are kept fixed. The natural axiom for this case, *population monotonicity*, requires that when new agents join the economy (for a fixed number of facilities), all the agents that were initially present weakly lose. This property is studied in this model by Miyagawa (1998), and by Ehlers (2003) in a model similar to this one. Both restrict attention to the case of two facilities. Gordon (2006) shows that this axiom has strong general implications in the broader class of pure public goods, of which this model is a special case.

all weakly gain from this double expansion. In fact, we will restrict attention to changes where exactly one facility is brought into the economy and exactly one peak-unanimous profile is added to the initial profile. Observe that the additional facility is indeed sufficient to fully satisfy the newcomers, since it could be located at their common peak.

A rule f satisfies **object-population monotonicity** if, for each problem (k, R_N) with $k < |p(R_N)|$, for each peak-unanimous profile $R_M \in \mathcal{T}$ such that $N \cap M = \emptyset$, we have $f_{k+1}(R_N, R_M) R_N f_k(R_N)$.

Although all the priority rules presented in Examples 1 to 6 satisfy this axiom, not all priority rules do, as shown in the following example.

Example 7: Let \succ be the priority such that for all two compatible peak-unanimous profiles R_L and R_M , we have $R_L \succ R_M$ if either $(|L| > |M|$ and $|L| \geq 3)$ or $(|L| = 1$ and $|M| = 2)$ or $(|L| = |M|$ and $p(R_L) < p(R_M))$. The priority rule f defined by \succ violates object-population monotonicity. To see this, consider a profile $R_{\{1,2,3,4\}}$ such that $p(R_1) = p(R_2) < p(R_3) < p(R_4)$. Therefore $R_1 \succ R_3 \succ R_4 \succ R_{\{1,2\}}$. Thus, $f_1(R_{\{1,3,4\}}) = p(R_1)$ and $f_2(R_{\{1,2,3,4\}}) = \{p(R_3), p(R_4)\}$. Therefore, $f_1(R_{\{1,3,4\}}) P_1 f_2(R_{\{1,2,3,4\}})$, in contradiction with object-population monotonicity.

More generally, one can show that a priority rule satisfies object-population monotonicity if and only if the priority \succ that defines it is such that, bringing more agents into a peak-unanimous profile cannot lower its priority by more than one step along any chain of compatible peak-unanimous profiles. In other words, there are no four peak-unanimous profiles R_M, R_K, R_H and R_L such that $p(R_M) = p(R_L)$ and $R_M \succ R_H \succ R_K \succ R_{M \cup L}$. In example 7, the violation occurs because the axiom requires that agent 1, who has the same peak as the newcomer agent 2, also weakly gain from the change. This suggests the following

weakening of the axiom. Consider any profile enlargement. Among the agents initially present, some may have the same peak location as the newcomers (call them allies), and others may not (call them opponents). The weaker version of the axiom requires that all opponents weakly gain, and requires nothing for allies.

A rule f satisfies **object-population monotonicity**₋ if, for each problem (k, R_N) with $k < |p(R_N)|$, for each peak-unanimous profile $R_M \in \mathcal{T}$ such that $N \cap M = \emptyset$, and for all $i \in N$ such that $p(R_i) \neq p(R_M)$, we have $f_{k+1}(R_N, R_M) R_N f_k(R_N)$.

It turns out that any priority rule satisfies this weaker version of the axiom.¹⁸ This observation brings out the following question: are all rules that satisfy efficiency and object-population monotonicity₋ priority rules? The answer is negative, as shown in the following example.

Example 8: Let \triangleright be the strict ordering on \mathbb{R} such that, for all $x, x', y, y', z, z' \in \mathbb{R}$ satisfying $x < x' < 0 < y < y' < 1 < z < z'$, we have $0 \triangleright 1 \triangleright y \triangleright y' \triangleright x \triangleright x' \triangleright z \triangleright z'$. Let \succ be the priority such that, for any two peak-unanimous profiles R_L and R'_M , we have $R_L \succ R'_M$ if and only if $p(R_L) \triangleright p(R'_M)$. Let g be the priority rule defined by \succ . Let f be the rule such that, for each problem (k, R_N) , if $k = 1$ and $p(R_N) = \{-1, 1\}$, then $f_k(R_N) \equiv \{\frac{1}{2}\}$, and otherwise $f_k(R_N) \equiv g_k(R_N)$. The rule f satisfies efficiency and object-population monotonicity₋, but it is not a priority rule.

¹⁸Even the strong version of the axiom does not imply object-monotonicity. However, object-population monotonicity₋, efficiency and sovereignty do imply object-monotonicity. Both versions of object-population monotonicity are also logically independent from population monotonicity. The rules presented in Example 5 satisfy object-population monotonicity, but not population-monotonicity. Moreover consider the rule f for which for each problem $(1, R_N)$, we have $f_1(R_N) = f_1^{LP}(R_N)$, and for each $k \geq 2$ and each problem (k, R_N) , we have $f_k(R_N) = f_k^{RP}(R_N)$. This rule satisfies population monotonicity but not object-population monotonicity₋.

We wish to eliminate rules that present pathologies, as in Example 8. This leads us to introduce a third axiom. This axiom only restricts the one-facility component f_1 of a rule, as the population varies. It says that the choice y made by the rule f_1 for a given profile R_N can be reversed in favor of any given location x provided that some appropriately selected peak-unanimous profile R_M is brought into the economy in support of x .

A rule f satisfies **sovereignty₋** if, for each profile R_N , each location $x \in \mathbb{R} \setminus f_1(R_N)$, there exists a peak-unanimous profile $R_M \in \mathcal{T}$ such that $M \cap N = \emptyset$, that satisfies $f_1(R_N, R_M) = \{x\} = p(R_M)$.

In fact, we will use the following slightly stronger version of the axiom, which requires that there are infinitely many distinct and compatible such peak-unanimous profiles.

A rule f satisfies **sovereignty** if, for each profile R_N , each location $x \in \mathbb{R} \setminus f_1(R_N)$, and each population L , there exists a peak-unanimous profile $R_M \in \mathcal{T}$ such that M is disjoint from both L and N , that satisfies $f_1(R_N, R_M) = \{x\} = p(R_M)$.

The rule presented in Example 8 violates both versions of the axiom. In fact, we will show in the next section that any efficient, object-population monotonic₋ and sovereign rule is a priority rule. However, not all priority rules are sovereign. While the rules in examples 5, 6 and 7, and some of the rules in examples 3 and 4 are, the ones in examples 1 and 2, and some of the rules in example 3 and 4 are not. We leave it to the reader to verify that a priority \succ defines a sovereign priority rule if and only if the following two conditions hold. (i) For all peak-unanimous $R_H, R_K \in \mathcal{T}$ such that $R_H \succ R_K$, and for any population L , there exists a peak-unanimous profile $R_M \in \mathcal{T}$ such that M is disjoint from K and L , and satisfies $p(R_M) = p(R_K)$, and $R_{K \cup M} \succ R_H$. (ii) For each $R_H \in \mathcal{T}$

, each $x \neq p(R_H)$ and each population L , there exists a peak-unanimous profile $R_M \in \mathcal{T}$ such that $M \cap L = \emptyset$, and satisfies $p(R_M) = x$, and $R_M \succ R_H$. Let $\mathbb{P}_{\mathcal{T}}^* \subset \mathbb{P}_{\mathcal{T}}$ be the set of priorities over \mathcal{T} that satisfy these condition. Let $\mathcal{F}_S \subset \mathcal{F}$ be the set of priority rules that are defined by a priority in $\mathbb{P}_{\mathcal{T}}^*$.

6 Main characterization

We are now ready to present our main result.

Theorem 1: *A rule f satisfies efficiency, object-population monotonicity₋ and sovereignty if and only if $f \in \mathcal{F}_S$.*

As we observed in Section 5, it is straightforward to verify that all the elements of the class \mathcal{F}_S satisfy the three axioms. Before proving the converse, we first present two useful lemmas. The first one states that *object-population monotonicity₋* and *sovereignty* imply the following property.

A rule f satisfies **strong sovereignty** if for each problem (k, R_N) , each location $x \in \mathbb{R} \setminus f_k(R_N)$, and each population L , there exists a peak-unanimous profile R_M such that M is disjoint from both L and N , that satisfies $p(R_M) = \{x\} \subseteq f_k(R_N, R_M)$.

Lemma 1: *If f satisfies object-population monotonicity₋ and sovereignty, then it satisfies strong sovereignty*

Proof. Let (k, R_N) be an arbitrary problem, let $x \in \mathbb{R}$ be an arbitrary location, and let L be an arbitrary population. We will prove that there exists a peak-unanimous profile R_M such that $M \cap (L \cup N) = \emptyset$ and $p(R_M) = \{x\} \subseteq f_k(R_N, R_M)$. If $k = 1$, then *sovereignty* guarantees that this is true. So suppose instead, that $k > 1$. Let $\ell \equiv |p(R_N)|$. We have $\ell \leq k$.

Let $R_{N_1}, \dots, R_{N_{\ell-(k-1)}}$ be $\ell - (k - 1)$ distinct maximal peak-unanimous sub-profiles of R_N . Let $\tilde{N} \equiv N_1 \cup \dots \cup N_{\ell-(k-1)}$. We have $|p(R_{\tilde{N}})| = \ell - (k - 1)$.

Let $\tilde{L} \equiv (N \setminus \tilde{N}) \cup L$. By *sovereignty*, there exists a profile $R_M \in \mathcal{R}^M$ such that $M \cap (\tilde{L} \cup \tilde{N}) = \emptyset$, and $p(R_M) = \{x\} = f_1(R_{\tilde{N}}, R_M)$. We have $|p(R_{N \setminus \tilde{N}})| = k - 1$. By *object-population monotonicity₋*, applied $k - 1$ times, we have in particular $f_k(R_{\tilde{N}}, R_M, R_{N \setminus \tilde{N}}) R_M f_1(R_{\tilde{N}}, R_M)$. Since $f_1(R_N, R_M) = \{x\} = p(R_M)$, this implies that $x \in f_k(R_{\tilde{N}}, R_M, R_{N \setminus \tilde{N}})$, i.e. $p(R_M) = \{x\} \subseteq f_k(R_N, R_M)$. Since $M \cap (L \cup N) = M \cap (\tilde{L} \cup \tilde{N}) = \emptyset$, therefore the population M satisfies all the desired properties. ■

The second lemma shows that the three axioms of the theorem imply that each public facility must be located at some agent's peak location.

A rule f satisfies **peak-selection** if for each problem (k, R_N) , we have $f_k(R_N) \subseteq p(R_N)$.

Lemma 2: *If f satisfies efficiency, object-population monotonicity₋ and sovereignty, then it satisfies peak-selection.*

Proof: Let f satisfy the three axioms. We prove by induction on k , that each f_k satisfies *peak-selection*.

Step 1. The rule f_1 satisfies peaks-selection.

Suppose by contradiction that f_1 does not satisfy *peak-selection*. Let $R_N \in \mathcal{R}^N$ and let $x \in \mathbb{R}$ be such that $f_1(R_N) = \{x\}$ and $x \notin p(R_N)$. Throughout the proof, for all $j \in N$, let $p_j := p(R_j)$. By *efficiency*, there are two agents $i, \ell \in N$ such that $p_i < x < p_\ell$, and $p(R_N)$ has no element strictly comprised between p_i and p_ℓ . Let $u, v \in \mathbb{R}$ be such that $p_i < u < x < v < p_\ell$. By *strong sovereignty*, there are peak-unanimous profiles R_J and R_K satisfying the following conditions. Let R_J be such that $J \cap N = \emptyset$, $p(R_J) = \{u\}$, and $u \in f_2(R_N, R_J)$. Similarly, let R_K be such that $K \cap (N \cup J) = \emptyset$, $p(R_K) = \{v\}$, and $v \in f_2(R_N, R_K)$. Let $A \equiv f_3(R_N, R_J, R_K)$. We will now show that A has at least four elements. First, by *object-population monotonicity₋*, $p(R_J) = \{u\}$ and $u \in$

$f_2(R_N, R_J)$ imply that $u \in A$. Similarly, by *object-population monotonicity*₋, $p(R_K) = \{v\}$ and $v \in f_2(R_N, R_K)$ imply that $v \in A$.

Second, by *object-population monotonicity*₋, since $f_1(R_N) = \{x\}$, then the set $f_2(R_N, R_K)$ has at least one location that is at least as good as x for R_i . Let $y \in f_2(R_N, R_K)$ such that $y R_i x$. *Efficiency* requires that $y \leq p_i$. In particular $y < v$. Since $|f_2(R_N, R_K)| = 2$, it then follows that $f_2(R_N, R_K) = \{y, v\}$. By *efficiency*, there exists an agent $h \in N$ such that $p_h \leq y$. By *object-population monotonicity*₋, since $f_2(R_N, R_K) = \{y, v\}$, then the set A has at least one location that is at least as good as y for R_h . Let $y' \in A$ such that $y' R_h y$. Since $p_h \leq y$, then in particular $y' \leq y < v$.

Third, by *object-population monotonicity*₋, since $f_1(R_N) = \{x\}$, then the set $f_2(R_N, R_J)$ has at least one location that is at least as good as x for R_ℓ . Let $z \in f_2(R_N, R_J)$ such that $z R_\ell x$. *Efficiency* requires that $p_\ell \leq z$. In particular $u < z$. Since $|f_2(R_N, R_J)| = 2$, it then follows that $f_2(R_N, R_J) = \{u, z\}$. By *efficiency*, there exists an agent $m \in N$ such that $z \leq p_m$. By *object-population monotonicity*₋, since $f_2(R_N, R_J) = \{u, z\}$, then the set A has at least one location that is at least as good as x for R_m . Let $z' \in A$ such that $z' R_m z$. Since $z \leq p_m$, then in particular $u < z \leq z'$.

In conclusion, $y' < u < v < z'$ are four distinct elements of A , in contradiction with $|A| = 3$.

Step 2. Let $k \geq 1$. Suppose that f_k satisfies peak-selection. Then f_{k+1} also satisfies peak-selection.

Let $R_N \in \mathcal{R}^N$. If $|p(R_N)| = k + 1$, then by *efficiency*, $f_{k+1}(R_N) = p(R_N)$, so the claim is true. Suppose then that $|p(R_N)| > k + 1$. We will first show that $f_{k+1}(R_N) \cap p(R_N) \neq \emptyset$. Since $|p(R_N)| > k + 1$, then $p(R_N) \not\subseteq f_{k+1}(R_N)$. Let R_M be a maximal peak-unanimous subprofile of R_N , with $M \subset N$, such that $p(R_M) \not\subseteq f_{k+1}(R_N)$. Consider the profile $R_{N \setminus M}$. We have $|p(R_{N \setminus M})| > k$. From the induction hypothesis, we have $f_k(R_{N \setminus M}) \subseteq p(R_{N \setminus M})$. By *object-*

population monotonicity₋, we have $f_k(R_{N \setminus M}) \subseteq f_{k+1}(R_N)$. Thus $f_k(R_{N \setminus M}) \subseteq f_{k+1}(R_N) \cap p(R_N)$. Therefore, $f_{k+1}(R_N) \cap p(R_N) \neq \emptyset$.

Let $x \in f_{k+1}(R_N) \cap p(R_N)$. Let R_L be the maximal peak-unanimous subprofile of R_N such that $p(R_L) = \{x\}$. Then $|p(R_{N \setminus L})| > k$ and $x \notin p(R_{N \setminus L})$. From the induction hypothesis, we have $f_k(R_{N \setminus L}) \subseteq p(R_{N \setminus L})$. By *object-population monotonicity*₋, we have $f_k(R_{N \setminus L}) \subseteq f_{k+1}(R_N)$. Since $x \in f_{k+1}(R_N)$, then $f_k(R_{N \setminus L}) \cup \{x\} \subseteq f_{k+1}(R_N)$. Since $x \notin p(R_{N \setminus L})$, then $x \notin f_k(R_{N \setminus L})$. Thus, $|f_k(R_{N \setminus L}) \cup \{x\}| = k+1 = |f_{k+1}(R_N)|$. Therefore, $f_{k+1}(R_N) = f_k(R_{N \setminus L}) \cup \{x\}$. Therefore, $f_{k+1}(R_N) \subseteq p(R_N)$, the desired conclusion. ■

We are now ready to prove Theorem 1.

Proof of Theorem 1 (only if implication): Let f be an arbitrary rule that satisfies *efficiency*, *object-population monotonicity*₋ and *sovereignty*.

Step 1: Construction of a candidate priority $\succ \in \mathbb{P}_{\mathcal{T}}^$ from f .*

By *peak-selection*, for each two compatible peak-unanimous profiles R_L and R_M , we have $f_1(R_L, R_M) \subset p(R_L) \cup p(R_M)$. Let \succ be the binary relation over peak-unanimous profiles such that, for each two compatible peak-unanimous profiles R_L and R_M , we have $R_L \succ R_M$ if $f_1(R_L, R_M) = p(R_L)$. By construction, the relation \succ is *asymmetric* and *almost complete*. It remains to show that \succ is *almost transitive*. Consider three arbitrary peak-unanimous profiles R_K , R_L and R_M such that R_K and R_M are compatible. Suppose that $R_K \succ R_L$ and $R_L \succ R_M$. Then in particular, R_K and R_L are compatible, and R_L and R_M are compatible. By definition of \succ , we know that $f_1(R_K, R_L) = p(R_K)$ and $f_1(R_L, R_M) = p(R_L)$. This and *object-population monotonicity*₋ imply that $p(R_K) \subset f_2(R_K, R_L, R_M)$ and $p(R_L) \subset f_2(R_K, R_L, R_M)$. By compatibility, we have $p(R_K) \neq p(R_L)$. Therefore $f_2(R_K, R_L, R_M) = p(R_K) \cup p(R_L)$. But this and *object-population monotonicity*₋ imply that $f_1(R_K, R_M) = p(R_K)$, i.e. $R_K \succ R_M$, the desired conclusion. Therefore \succ is a priority in $\mathbb{P}_{\mathcal{T}}$.

It only remains to show that \succ belongs to the subset \mathbb{P}_T^* . Let R_H and R_K be two arbitrary peak-unanimous profiles satisfying $R_H \succ R_K$, and let L be an arbitrary population. By *sovereignty*, since $f_1(R_H, R_K) \neq p(R_K)$, then there exists a peak-unanimous profile R_M satisfying $M \cap (H \cup K \cup L) = \emptyset$, and $f_1(R_H, R_K, R_M) = p(R_M) = p(R_K)$. Since (R_K, R_M) is then a peak-unanimous profile, we have $(R_K, R_M) \succ R_H$. Therefore $\succ \in \mathbb{P}_T^*$.

Step 2: Let $\succ \in \mathbb{P}_T^$ be defined as in step 1 from f . Let $R_N \in \mathcal{P}_1$. Let R_M and R_L be distinct maximal peak-unanimous subprofiles of R_N such that $R_M \succ R_L$. Then $f_1(R_N) \neq p(R_L)$.*

The proof is by induction on $|p(R_N)|$. For $|p(R_N)| = 2$, the claim follows from the definition of \succ . Let $R_N \in \mathcal{P}_1$ be such that $|p(R_N)| = n \geq 3$, and suppose that the claim is true for all $R'_{N'} \in \mathcal{P}_1$ such that $|p(R'_{N'})| = n - 1$. Since $n \geq 3$, there is a maximal peak-unanimous subprofile R_H of R_N distinct from both R_L and R_M , such that $f_1(R_N) \neq p(R_H)$. By the induction hypothesis, $f_1(R_{N \setminus H}) \neq p(R_L)$. Let R_K be the maximal peak-unanimous subprofile of R_N , distinct from both R_L and R_H , such that $f_1(R_{N \setminus H}) = p(R_K)$. *Strong sovereignty* ensures that there exists a peak-unanimous profile R_J such that $N \cap J = \emptyset$, $p(R_J) = p(R_H)$, and $p(R_J) \subset f_2(R_N, R_J)$. By *object-population monotonicity₋*, we have $p(R_K) \subset f_2(R_N, R_J)$. Since $p(R_K) \neq p(R_J)$, thus $f_2(R_N, R_J) = p(R_K) \cup p(R_J)$. By *object-population monotonicity₋*, we have $f_1(R_N) \subset p(R_K) \cup p(R_H)$. Therefore, $f_1(R_N) \neq p(R_L)$, i.e. the claim is true for $|p(R_N)| = n$. Thus, it is true for any integer value of $|p(R_N)|$, which proves the claim.

Step 3: Let \succ be defined as in step 1 from f . Let (k, R_N) be a problem. Let R_M and R_L be distinct maximal peak-unanimous subprofiles of R_N such that $R_M \succ R_L$ and $p(R_L) \subseteq f_k(R_N)$. Then $p(R_M) \subseteq f_k(R_N)$.

Let R_M and R_L satisfying the assumptions of the step. Consider all the max-

imal peak-unanimous subprofiles of R_N that are distinct from R_M and whose peak location is not contained in $f_k(R_N)$. There are exactly $m := |p(R_N)| - k - 1$ such subprofiles. Let R_{M_1}, \dots, R_{M_m} be these subprofiles. By step 2, we know that $f_1(R_{M_1}, \dots, R_{M_m}, R_M, R_L) \neq p(R_L)$. Therefore $f_1(R_{M_1}, \dots, R_{M_m}, R_M, R_L) \subset p(R_{M_1 \cup \dots \cup M_m \cup M})$. Next, using *object-population monotonicity* $k - 1$ times, we obtain that $f_k(R_N) \cap p(R_{M_1 \cup \dots \cup M_m \cup M}) \neq \emptyset$. Since $f_k(R_N) \cap p(R_{M_1 \cup \dots \cup M_m}) = \emptyset$ by definition of the subprofiles R_{M_k} , therefore $p(R_M) \subseteq f_k(R_N)$, the desired conclusion.

It follows from Step 3 that any rule satisfying the axioms is the priority rule associated with \succ , where \succ is defined as in Step 1. ■

Independence of the axioms We verify that the axioms are independent. First, the left-peaks rule f^{LP} (like any other rule $\mathcal{F} \setminus \mathcal{F}_S$) satisfies all the axioms but *sovereignty*. Second, the rule f such that f_1 selects the location x_m of the (left)-median voter when $k = 1$, and such that f_k coincides with the left-peaks rule for all $k \geq 2$ satisfies *efficiency* and *sovereignty*, but violates *object-population monotonicity*. Finally, the rule that locates the first facility on the left-majoritarian peak and the remaining facilities at the smallest $k - 1$ positive integers distinct from $p(R_N)$ satisfies *object-population monotonicity* and *sovereignty* but violates *efficiency*.¹⁹

7 Sovereign and non-sovereign priority rules

Theorem 1 characterizes a proper subset \mathcal{F}_S of the set \mathcal{F} of all priority rules. We already pointed out that all rules in $\mathcal{F} \setminus \mathcal{F}_S$ satisfy *efficiency* and *object-population monotonicity*, but violate *sovereignty*. The motivation for introduc-

¹⁹Observe that the three axioms imply object-monotonicity, since all priority rules are object-monotonic. Also, none of the rules in \mathcal{F}_S satisfies population-monotonicity, since the only population-monotonic priority rules are the left peaks rule and the right peaks rule and these two rules are not in \mathcal{F}_S .

ing sovereignty was to eliminate pathological rules that also satisfy *efficiency* and *object-population monotonicity₋*, but are not priority rules, such as the rule in Example 8. Requiring *sovereignty* eliminates all such rules, but also eliminates some priority rules, the ones in $\mathcal{F} \setminus \mathcal{F}_S$. In this section, we show that the violation of sovereignty is, in some sense, more severe for rules that are not priority rules than for non-sovereign priority rules.

The domain of any rule is the set $\bigcup_{k=1}^{\infty} [\{k\} \times \mathcal{P}_k]$. We say that a rule f *finitely satisfies* a property Z if, for all finite subset D of $\bigcup_{k=1}^{\infty} [\{k\} \times \mathcal{P}_k]$, there is a rule f^D that satisfies Z and coincides with f on D . Obviously, any rule that satisfies Z also finitely satisfies Z . Intuitively, if a rule f does not satisfy an axiom, but does satisfy it finitely, one should consider that the violation is less severe, than if it does not even satisfy it finitely.

The gap between satisfying and finitely satisfying a property depends on the property itself. For some axioms, there is none. For example, if a rule finitely satisfies efficiency, then it is efficient. If a rule finitely satisfies object-population monotonicity₋, then it is object-population monotonic₋. In contrast, it can be shown that all rules finitely satisfy sovereignty,²⁰ while not all rules are sovereign.

We are interested in rules that finitely satisfy the combination of our three axioms. These are the rules that coincide on any finite subset D of $\bigcup_{k=1}^{\infty} [\{k\} \times \mathcal{P}_k]$ with some rule that satisfies all three axioms,²¹ i.e. some rule in \mathcal{F}_S . We have the following result.

Theorem 2: *A rule f is an element of \mathcal{F} if and only if it finitely satisfies the combination of efficiency, object-population monotonicity₋ and sovereignty.*

²⁰The proof of this claim is similar to the proof of Theorem 2.

²¹This is stronger than the requirement that, for each of the three axioms, the rule should coincide on any finite subset D of $\bigcup_{k=1}^{\infty} [\{k\} \times \mathcal{P}_k]$ with some rule that satisfies this axiom alone. This weaker requirement is equivalent to efficiency and object-population monotonicity₋.

Proof: if implication. Let f be a rule that finitely satisfies the combination. Then f coincides on any finite subset D of $\bigcup_{k=1}^{\infty} [\{k\} \times \mathcal{P}_k]$ with some priority rule in \mathcal{F}_S . Thus, for each problem (k, R_N) , $f_k(R_N)$ is the image of R_N by the k -th component of some priority rule in \mathcal{F}_S . Thus, in particular, $f_k(R_N) \subseteq p(R_N)$. Therefore, f satisfies *peaks-selection*. Next, for all two problems (k, R_N) and $(k', R'_{N'})$, since f coincides with some priority rule on $D = \{(k, R_N), (k', R'_{N'})\}$, then there is no pair of locations $(x, y) \in \mathbb{R}^2$ such that $x \in f_k(R_N) \setminus f_{k'}(R'_{N'})$ and $y \in f_{k'}(R'_{N'}) \setminus f_k(R_N)$. Thus, f is a priority rule.

Only if implication. Let f be any arbitrary element of \mathcal{F} . Let \succ be the priority associated with f . Let D be an arbitrary finite subset of $\bigcup_{k=1}^{\infty} [\{k\} \times \mathcal{P}_k]$. We will construct a priority rule $g \in \mathcal{F}_S$, such that f and g coincide on D . Let $I_D := \{i \in \mathbb{N} : i \in N, \text{ for some } (k, R_N) \in D\}$. Let \succ' be the priority such that for all compatible R_M and R_N , we have $R_M \succ' R_N$ if and only if either $(N \cup M \subseteq I_D \text{ and } R_M \succ R_N)$ or $(\exists i \in N \setminus I_D \text{ such that } \forall j \in M \setminus I_D, \text{ we have } i > j)$. Clearly f and g coincide on D , and $g \in \mathcal{F}_S$, the desired conclusion. ■

Theorem 2 has an interesting topological interpretation. For each $k \in \mathbb{N}$, let us endow \mathcal{X}_k with the discrete topology (all subsets of \mathcal{X}_k are open). Furthermore, let us endow $\prod_{k=1}^{\infty} \mathcal{X}_k^{\mathcal{P}_k}$, the set of all rules, with the corresponding product topology, which we label the product-discrete topology. Then a rule finitely satisfies a property Z if and only if f belongs to the product-discrete closure of the set of rules satisfying Z . Theorem 2 says that \mathcal{F} is the product-discrete closure of \mathcal{F}_S . Moreover, we pointed out that the set of all rules is the closure of the set of sovereign rules. As for the set of efficient rules, the set of object-population monotonic_ rules, and the set of priority rules, they all are closed.

8 Other properties

In this section and the next one, we investigate on subclasses of priority rules in \mathcal{F} that satisfy other additional desirable properties, such as *anonymity*, *hiding-proofness* and *median-selection*. In the next section, we will turn our attention to *strategy-proofness*.

A rule f satisfies **anonymity** if, for all $k \geq 1$, and all $R_N, R'_M \in \mathcal{P}_k$ where for all $R \in \mathcal{R}$, we have $|\{i \in M : R'_i = R\}| = |\{i \in N : R_i = R\}|$, we have $f_k(R_N) = f_k(R'_M)$. A priority \succ is *anonymous* if it satisfies the following condition. For all $R_M, R_N, R'_{M'}, R'_{N'} \in \mathcal{T}$, such that: (i) R_M and R_N are compatible, (ii) $R'_{M'}$ and $R'_{N'}$ are compatible, (iii) for all $R \in \mathcal{R}$, we have $|\{i \in M : R_i = R\}| = |\{i \in M' : R'_i = R\}|$, (iv) for all $R \in \mathcal{R}$, we have $|\{i \in N : R_i = R\}| = |\{i \in N' : R'_i = R\}|$, the following equivalence holds

$$R_M \succ R_N \Leftrightarrow R'_{M'} \succ R'_{N'}.$$

We leave it to the reader to verify the following result.

Proposition 1: *A priority rule $f \in \mathcal{F}$ satisfies anonymity if and only if its priority \succ is anonymous.*

Next, we consider a property that ensures that no agent gains from hiding from the social planner. Consider a university which has to choose k software licenses from a larger set of available software for its employees. The decision is taken through an online poll, to which employees can freely respond. One would want that no employee has an incentive not to participate in the poll. A rule f satisfies **hiding-proofness** if for each problem $k \geq 1$, each $R_N \in \mathcal{P}_k$, and each $i \in N$ such that $R_{N \setminus \{i\}} \in \mathcal{P}_k$, we have $f_k(R_N) R_i f_k(R_{N \setminus \{i\}})$. A priority \succ is *monotone in population* if, for all peak-unanimous profiles R_L, R_N such that R_N and R_L are compatible, and any agent $i \in N$, we have $R_{N \setminus \{i\}} \succ R_L \Rightarrow R_N \succ R_L$. The following holds.

Proposition 2: *A priority rule $f \in \mathcal{F}$ satisfies hiding-proofness if and only if its priority \succ is monotone in population.*

Proof. The *if* implication is clear. Let us prove the *only if* implication. Let f be a priority rule that satisfies *hiding-proofness*, associated with the priority \succ . Let $R_L, R_N \in \mathcal{T}$ be such that R_N and R_L are compatible. Let $i \in N$. Suppose that $R_{N \setminus \{i\}} \succ R_L$. Then $f_1(R_L, R_{N \setminus \{i\}}) = p(R_{N \setminus \{i\}}) = p(R_N)$. This equality and *hiding-proofness* imply $f_1(R_L, R_N) = p(R_N)$, which in turn implies $R_N \succ R_M$, the desired conclusion. ■

Last, an important rule for selecting a single location in this model is the one which selects the median peak, for each preference profile R_N such that $|N|$ is odd (Black, 1948). A natural question is thus whether there exists a priority rule f such that, for any R_N such that $|N|$ is odd, the location $f_1(R_N)$ is the median of the peaks. Unfortunately, the answer is negative. In fact we are able to show a much stronger negative result.

Proposition 3: *Let $N = \{1, \dots, 11\}$. There is no rule f satisfying object-population monotonicity₋ and such that, for each $R_M \in \mathcal{R}^M$, with $M \subset N$ such that $|M|$ is odd, the location $f_1(R_M)$ is the median of the peaks of the profile R_M .*

Proof: By contradiction, suppose that f is such a rule. Consider a profile R_N whose only maximal peak-unanimous subprofiles $R_{\{1,2,3,4\}}$, $R_{\{5,6\}}$, $R_{\{8,9,10,11\}}$ have peaks locations respectively at 1, 2 and 4. Let R_7 be such that $p(R_7) = 3$. Then, by the median property, we have $f_1(R_{N \setminus \{1,2,3,4\}}) = 4$, $f_1(R_{N \setminus \{5,6\}}) = 3$ and $f_1(R_{N \setminus \{8,9,10,11\}}) = 1$. By object-population monotonicity₋, we have $\{1, 3, 4\} \subseteq f_2(R_N)$, which contradicts $|f_2(R_N)| = 2$. ■

9 Strategy-proofness and hierarchical rules

In this section, we investigate on the class of procedures to locate a single facility that are both the first component of some priority rule and robust to the manipulation of preferences. A 1-rule is a mapping $g : \mathcal{P}_1 \rightarrow \mathcal{X}_1$. It can also be seen as the first component of a rule. A 1-rule is a 1-priority rule if it is the first component of a priority rule. The non-manipulability property we are interested in requires that truthfully reporting preferences be a (weakly) dominant strategy in the associated direct revelation game:

A 1-rule g satisfies **strategy-proofness** if for each $R_N \in \mathcal{P}_1$, each $i \in N$ and each $R'_i \in \mathcal{R}$, we have $g(R_N) R_i g(R'_i, R_{-i})$.

The first components of the left-peaks rule, the right-peaks rule and of all the serial dictatorships are strategy-proof priority 1-rules. More generally, the first component of any hierarchical rule (see Example 8) is strategy-proof. In fact, we will show in this section that the hierarchical 1-rules are the only strategy-proof priority 1-rules.²²

9.1 Strategy-proof priority N -schemes

For each population N , an N -scheme is a mapping $g_N : \mathcal{R}^N \rightarrow \mathcal{X}_1$. It can also be seen as the restriction of a 1-rule to the subset \mathcal{R}^N . Moulin (1980) introduced the following class of N -schemes.

An N -scheme g_N is a *peak-selecting generalized median* N -scheme if there is a list $\{a_M\}_{M \subseteq N}$ of parameters such that $a_N = -\infty$, $a_\emptyset = +\infty$, and for each $M \subseteq N$, we have $a_M \in \{-\infty, +\infty\}$; for all $L, M \subseteq N$, we have $L \subseteq M \Rightarrow a_M \leq$

²²The definition of strategy-proofness can be extended to rules in the obvious way. See the comments at the end of the section.

a_L ; and for each $R_N \in \mathcal{R}^N$, we have

$$g(R_N) = \min_{M \subseteq N} \left\{ \max_{i \in M} [a_M, p(R_i)] \right\}.$$

It is well known that peak-selecting generalized median N -schemes are the only peak selecting strategy-proof N -schemes.²³ Let \mathcal{M}_N be the class of peak-selecting generalized median N -schemes. We say that a subset $M \subseteq N$ is *minimal* for the N -scheme g if $a_{M,N} = -\infty$ and, for all $L \subsetneq M$, we have $a_L = +\infty$. In particular, any minimal subset M is nonempty, and is therefore a subpopulation of N . Clearly, we have

$$g(R_N) = \min_{\substack{M \subseteq N, \\ \text{s.t. } M \text{ is minimal}}} \{ \max [p(R_M)] \}.$$

We say that an N -scheme g_N is a *priority N -scheme* if there is a priority 1-rule that coincides with g_N on \mathcal{R}^N . We now provide a description of all strategy-proof priority N -schemes, in terms of the collection of their minimal sets.

Lemma 3: *Let g_N be a strategy-proof priority N -scheme. Then, g_N is peak-selecting generalized median. Furthermore, either g_N has exactly one minimal set, or all minimal sets of g_N are singletons.*

Proof: Since g_N is a priority N -scheme, it satisfies peaks-selection. Since it is also strategy-proof, then it is a peak-selecting generalized median. Therefore, it can be described by its minimal sets. By contradiction, suppose that L and M are two distinct minimal sets for g_N , and that $|M| \geq 2$. Let $j^* \in M \setminus L$. Since $|M| \geq 2$, then $M \setminus \{j^*\} \neq \emptyset$. Since L and M are two distinct minimal sets for g_N , then $L \subsetneq M$, which further implies that $N \setminus M \neq \emptyset$. The sets $\{j^*\}$, $M \setminus \{j^*\}$, and $N \setminus M$ form a disjoint partition of N , and none of these sets is empty.

²³This follows from the work of Moulin (1980) and Barberà and Jackson (1994). This family plays an important role in the work of Jackson and Nicolò (2004), and Cantala (2004).

Let $R_N \in \mathcal{R}^N$ be a profile whose maximal peak-unanimous subprofiles are $R_{M \setminus \{j^*\}}$, $R_{N \setminus M}$ and R_{j^*} and have peak locations respectively at 0, 1 and 2, so that $p(R_N) = \{0, 1, 2\}$. Since M is minimal, then for each minimal set K , we do not have $K \subsetneq M$. In particular, we do not have $K \subseteq M \setminus \{j^*\}$. Therefore $\max[p(R_K)] \geq 1$. Since this holds for each minimal set K , therefore $g_N(R_N) \geq 1$. Since, however, $\max[p(R_L)] = 1$ and $L \subseteq N \setminus \{j\}$, then $g_N(R_N) = 1$. Thus, $R_{N \setminus M} \succ R_{M \setminus \{j^*\}}$.

Next, let R'_{j^*} be a preference whose peak location is at -1 . Let $R'_N := (R'_{j^*}, R_{N \setminus \{j^*\}})$. Then, we have $p(R'_N) = \{-1, 0, 1\}$. Since M is minimal, then for each minimal set K , we do not have $K \subsetneq M$. In particular, we do not have $K \subseteq \{j^*\}$. Therefore $\max[p(R'_K)] \geq 0$. Since this holds for each minimal set K , therefore $g_N(R'_N) \geq 0$. Since, however, $\max[p(R'_M)] = 0$, then $g_N(R'_N) = 0$. Therefore, $R'_{M \setminus \{j^*\}} \succ R'_{N \setminus M}$, i.e. $R_{M \setminus \{j^*\}} \succ R_{N \setminus M}$, which contradicts $R_{N \setminus M} \succ R_{M \setminus \{j^*\}}$. ■

We now provide a convenient description of the family of strategy-proof priority N -schemes. Consider a pair (C, D) such that $C \subseteq N$ and $D \subseteq N$. Let the N -scheme g_N be the (C, D) -*oligarchy* if exactly one of the following mutually exclusive conditions holds.

- (a). $|C| \geq 2$, $D = \emptyset$, and for all $R^N \in \mathcal{R}^N$, $g_N(R_N) = \min\{p(R_C)\}$.
- (b). $C = \emptyset$, $|D| \geq 2$, and for all $R^N \in \mathcal{R}^N$, $g_N(R_N) = \max\{p(R_D)\}$.
- (c). $C = D = \{i^*\}$ for some $i^* \in N$, and for all $R^N \in \mathcal{R}^N$, $g_N(R_N) = p(R_{i^*})$.

The (C, D) -oligarchy is not defined for all pairs (C, D) , only the ones satisfying one of the above admissibility conditions. Also, two distinct pairs (C, D) and (C', D') define different oligarchies. In case (a), the scheme is similar to a left-peaks N -scheme, but it only takes into account a subset C of the population. In case (b), it is similar to a right-peaks N -scheme, but it takes into

account only a subset D of the population. In case (c), it is the dictatorship of some agent i^* . As we will see in Theorem 3, an N -scheme is a strategy-proof priority N -scheme if and only if it is a (C, D) -*oligarchy* for some admissible pair (C, D) .

9.2 Conjugate weak partial orders

In order to be able to characterize strategy-proof priority 1-rules, we introduce here a few concepts. Two weak partial orders \succeq_L and \succeq_R on \mathbb{N} are *conjugate* if for all $i, j \in \mathbb{N}$ such that $i \neq j$, we have $(i \succeq_L j) \Leftrightarrow \neg(j \succeq_R i)$. Given two weak partial orders \succeq' and \succeq'' on \mathbb{N} , let the *union* of \succeq' and \succeq'' be the binary relation \succeq on \mathbb{N} such that, for all $i, j \in \mathbb{N}$, we have $i \succeq j \equiv (i \succeq' j \text{ or } i \succeq'' j)$. The following holds.

Lemma 4: *Let \succeq_L and \succeq_R be two conjugate partial orders on \mathbb{N} . Then the union \succeq of \succeq_L and \succeq_R is a weak ordering.*

Proof. We leave it to the reader to verify that the relation \succeq is reflexive and complete. We will prove that it is transitive. Let i, j, k be arbitrary elements of \mathbb{N} , such that $i \succeq j$, and $j \succeq k$. We will prove that $i \succeq k$. If i, j and k are not three distinct integers, the claim is obviously true. Suppose then that i, j and k are three distinct integers. Four cases may arise. *Case 1:* $i \succeq_L j$, and $j \succeq_L k$. Then by transitivity of \succeq_L , we have $i \succeq_L k$. *Case 2:* $i \succeq_R j$, and $j \succeq_R k$. Then by transitivity of \succeq_R , we have $i \succeq_R k$. *Case 3:* $i \succeq_L j$, and $j \succeq_R k$. Since $i \succeq_L j$, $i \neq j$, and \succeq_L and \succeq_R are conjugates, therefore $\neg(j \succeq_R i)$. By transitivity of \succeq_R , and since $j \succeq_R k$ and $\neg(j \succeq_R i)$, therefore $\neg(k \succeq_R i)$. Since $k \neq i$, and \succeq_L and \succeq_R are conjugates, this is equivalent to $i \succeq_L k$. *Case 4:* $i \succeq_R j$, and $j \succeq_L k$. Following a reasoning symmetric to Case 3, we can prove that $i \succeq_R k$. In all four cases we proved, either that $i \succeq_L k$, or that $i \succeq_R k$. Therefore, in all four cases, $i \succeq k$, the desired conclusion, holds. ■

Lemma 5: *Let \succeq_L and \succeq_R be two conjugate weak partial orders. Let $T \subset \mathbb{N}$ be an indifference class for the union \succeq of \succeq_L and \succeq_R with at least two elements. Then, T is an indifference class, either for \succeq_L , or for \succeq_R , but not for both.*

Proof. Throughout the proof, let \bowtie , \bowtie_L , and \bowtie_R denote the respective indifference relations associated with \succeq , \succeq_L and \succeq_R . We will first show that for all $i, j \in T$, such that $i \neq j$, and $i \bowtie j$, we have either $i \bowtie_L j$ or $i \bowtie_R j$, but not both. Let $i, j \in T$ be two such agents. Since \succeq_L and \succeq_R are conjugates and $i \neq j$, we cannot have $i \succeq_L j$ and $j \succeq_R i$. For the same reason, we cannot have $i \succeq_R j$ and $j \succeq_L i$. Therefore, either $i \bowtie_L j$, or $i \bowtie_R j$, but not both. If T has exactly two elements, there is nothing left to prove. Assume then that T has at least three elements.

Next, we will show that for all three distinct $i, j, k \in T$, we have either $i \bowtie_L j \bowtie_L k$, or $i \bowtie_R j \bowtie_R k$, but not both. We will first show that $i \bowtie_L j \bowtie_R k$ cannot hold. Suppose, by contradiction, that these relations hold. Then, on the one hand, we have $i \succeq_L j \succeq_R k$. By Case 3 in the proof of Lemma 4, this implies that $i \succeq_L k$. But, also, on the other hand, we have $k \succeq_R j \succeq_L i$. By Case 4 in the proof of Lemma 4, this implies that $k \succeq_R i$. Thus we have $i \neq k$, $i \succeq_L k$ and $k \succeq_R i$, which contradicts that \succeq_L and \succeq_R are conjugates. Therefore, $i \bowtie_L j \bowtie_R k$ cannot hold. A symmetric reasoning shows that $i \bowtie_R j \bowtie_L k$ cannot hold either. Therefore, we have either $i \bowtie_L j \bowtie_L k$, or $i \bowtie_R j \bowtie_R k$, but not both. Since this holds for all distinct $i, j, k \in T$, it follows that T is an indifference class, either for \succeq_L , or for \succeq_R , but not for both, the desired conclusion. ■

9.3 Strategy-proof priority 1-rules

We are now ready to provide a characterization of all strategy-proof priority 1-rules. We describe this set in two ways.

Under the first description, each such 1-rule is viewed as the collection of

the N -schemes it induces: Each of them is a (C_N, D_N) -oligarchy. Moreover the parameters (C_N, D_N) are linked across population, in a way which will be made clear. We describe the exact nature of the linkage. In the second description, we focus on the priority itself and show that it is hierarchical, as defined in Example 4. We are now ready to present the main result of this section.

Theorem 3: *Let g be a 1-rule. The three following statements are equivalent.*

- (i). *g is a strategy-proof priority 1-rule.*
- (ii). *There are conjugate weak partial orders \succeq_L and \succeq_R such that, for each population N , the 1-rule g induces the (C_N, D_N) -oligarchic N -scheme g_N on N , with $C_N = \{i \in N : \forall j \in N, i \succeq_L j\}$ and $D_N = \{i \in N : \forall j \in N, i \succeq_R j\}$.*
- (iii). *g is a hierarchical 1-rule.*

Proof: (iii) \implies (i). A hierarchical rule g is obviously a priority 1-rule. It remains to show that it is strategy-proof. Let N be an arbitrary population. Let g_N be the N -scheme induced by g on N . Let T be the highest priority class of agents in the hierarchy such that $T \cap N \neq \emptyset$. Then either, for all profile $R_N \in \mathcal{R}^N$, we have $g_N(R_N) = \min [p(R_{N \cap T})]$ or, for all profile $R_N \in \mathcal{R}^N$, we have $g_N(R_N) = \max [p(R_{N \cap T})]$. In both cases, the N -scheme g_N is strategy-proof, the desired conclusion.

(i) \implies (ii). By Lemma 3, either g_N has a single minimal set and this set is not a singleton, or it contains more than one minimal set and they are all singletons, or it contains a single set and it is a singleton. Clearly, these three cases are mutually exclusive. In the first case, let C_N be the union of the many singleton sets for g_N , and let $D_N \equiv \emptyset$. In the second case, let D_N be unique (non singleton) minimal set, and let $C_N \equiv \emptyset$. In the third and last case, let i^* be the

unique element of the unique minimal set, and let $C_N \equiv \{i^*\}$, and $D_N \equiv \{i^*\}$. Clearly, for each N , the pair (C_N, D_N) satisfies conditions (a), (b) and (c) of the definition of an oligarchic 1-rule. Next, let \succeq_L and \succeq_R be two binary relations on \mathbb{N} defined as follows. For all $i, j \in \mathbb{N}$, let $(i \succeq_L j) \Leftrightarrow (i \in C_{\{i,j\}} \text{ or } i = j)$, and let $(i \succeq_R j) \Leftrightarrow (i \in D_{\{i,j\}} \text{ or } i = j)$. We will prove that \succeq_L, \succeq_R satisfy the requirements of statement (ii) in three steps.

Step 1. The binary relations \succeq_L and \succeq_R are transitive. We will only prove that \succeq_L is transitive. The proof that \succeq_R is transitive follows a symmetric reasoning. Let i, j and k be three agents such that $i \succeq_L j$ and $j \succeq_L k$. We claim that $i \succeq_L k$ holds. If these agents are not all distinct, then our claim is obvious. Suppose then that they are distinct. Let the preferences R_i, R_j and R_k have their respective peak locations at $-1, 0$ and 1 . Then, on the one hand, $i \succeq_L j$ and $i \neq j$ imply that $i \in C_{\{i,j\}}$. Since $p(R_i) = -1 < 0 = p(R_j)$, it follows that $g(R_i, R_j) = -1$. Since g is the priority rule associated with \succ , this implies that $R_i \succ R_j$. By a similar reasoning, we can show that $j \succeq_L k$ implies that $R_j \succ R_k$. By almost transitivity, we obtain $R_i \succ R_k$. Therefore, $g(R_i, R_k) = -1$. Since $p(R_i) = -1 < 1 = p(R_k)$, this implies that $i \in C_{\{i,k\}}$. Therefore, we have $i \succeq_L k$. Therefore, \succeq_L is transitive, the desired conclusion.

Step 2. The weak partial orders \succeq_L and \succeq_R are conjugate. Let $i, j \in N$ such that $i \neq j$ and $i \succeq_L j$. Then $i \in C_{\{i,j\}}$ and $i \neq j$. Therefore $D_{\{i,j\}} \subseteq \{i\}$. In particular, $j \notin D_{\{i,j\}}$ and $i \neq j$. Therefore, $\neg(j \succeq_R i)$, the desired conclusion.

Step 3. For all N , we have $C_N = \{i \in N : \forall j \in N, i \succeq_L j\}$ and $D_N = \{i \in N : \forall j \in N, i \succeq_R j\}$. We will only prove the claim for C_N . The proof of the claim for D_N follows a symmetric reasoning. First we prove that $C_N \subseteq \{i \in N : \forall j \in N, i \succeq_L j\}$. Let $i \in C_N$ and $j \in N$. We will prove that $i \succeq_L j$. Let R_N be a profile whose only maximal peak-unanimous subprofiles are R_i, R_j and $R_{N \setminus \{i,j\}}$ and have their respective peak locations at $-1, 0$ and 1 . Then $g(R_N) = \min[p(R_{C_N})] = -1$. This implies that $R_i \succ R_j$. Therefore, $g(R_{i,j}) =$

-1 , i.e. $i \succeq_L j$. Second, we prove that $\{i \in N : \forall j \in N, i \succeq_L j\} \subseteq C_N$. Let $i \in \{i \in N : \forall j \in N, i \succeq_L j\}$. We will prove that $i \in C_N$. Let j_1, \dots, j_l be the elements of $N \setminus i$. Let R_N be a profile such that $p(R_i) := 0$ and for all $k = 1, \dots, l$, we have $p(R_{j_k}) := l$. For all $k = 1, \dots, l$, we have $i \succeq_L j_k$, therefore $g(R_{i,j_k}) = -1$, therefore $R_i \succ R_{j_k}$. Since this holds for all $k = 1, \dots, l$ and g is a priority rule, it follows that $g(R_N) = -1 = p(R_i)$. Therefore $i \in C_N$, the desired conclusion.

(ii) \implies (iii). Let the statement (ii) hold. We will proceed in two steps. First, we will construct a candidate hierarchical priority \succ . Second, we will show that g is indeed the 1-rule associated with this priority.

Step 1: Definition of a candidate hierarchical priority. Let \succ be binary relation on \mathcal{T} defined as follows. Let \succeq be the union of \succeq_L and \succeq_R . Then \succeq is a weak ordering of \mathbb{N} , such that each indifference class W for \succeq is either an indifference class for \succeq_L , or for \succeq_R , but not for both. Let \succeq be the weak ordering on \mathcal{T} defined as follows. For all $R_M, R_N \in \mathcal{T}$, we have $R_M \succeq R_N$ if there is $i \in M$ such that for all $j \in N \cup M$, we have $i \succeq j$. Let \mathbb{T}^* be the set of non-singleton indifference classes for \succeq in \mathcal{T} . Then, for each non singleton indifference class $T \in \mathbb{T}^*$, there is a (unique) indifference class W_T for \succeq in \mathbb{N} , such that $T = \{R_N \in \mathcal{T} : N \cap W_T \neq \emptyset\}$. We know that for each class $T \in \mathbb{T}^*$, the class W_T satisfies exactly one of the following two mutually exclusives conditions. Either it is a non-singleton indifference class for \succeq_L , or it is a non-singleton indifference class for \succeq_R . We will now define a priority \succ_T over each non-singleton indifference class $T \in \mathbb{T}^*$. If W_T is a non-singleton indifference class for \succeq_L , then let \succ_T be the left peaks priority over T . If instead W_T is a non-singleton indifference class for \succeq_R , then let \succ_T be the right peaks priority over T . The tuple $(\succeq, (\succ_T)_{T \in \mathbb{T}^*})$ defines a hierarchical priority \succ .

Step 2: g is the priority 1-rule associated with \succ . We know that each of the N -schemes g_N satisfies peaks selection. Thus g satisfies peaks selection. Let

R_N be an arbitrary profile. Let R_L be the maximal peak-unanimous subprofile of R_N such that $g(R_N) = p(R_L)$. We will show that this implies that for all other maximal peak-unanimous subprofile R_M of R_N , we have $R_L \succ R_M$. To do so, we will distinguish three mutually exclusive cases. *Case 1:* $|C_N| \geq 2$. Then $p(R_L) = \min \{p(R_{C_N})\}$. Since it is nonempty, the set C_N is the highest indifference class for \succeq in N and is also an indifference class for \succeq_L . This and $p(R_L) = \min \{p(R_{C_N})\}$ implies that $R_L \succ R_M$, for all maximal peak-unanimous profiles R_M in R_N that are distinct from R_L . *Case 2:* $|D_N| \geq 2$. Then $p(R_L) = \max \{p(R_{D_N})\}$. Since it is nonempty, the set D_N is the highest indifference class for \succeq in N and is also an indifference class for \succeq_R . This and $p(R_L) = \max \{p(R_{D_N})\}$ implies that $R_L \succ R_M$, for all maximal peak-unanimous profiles R_M in R_N that are distinct from R_L . *Case 3:* $C_N = D_N = \{i\}$ for some $i \in N$. Then i is the sole element in the highest indifference class for \succeq in N . As a consequence, whichever maximal peak-unanimous profile of R_N has agent i enjoys the highest priority for \succ . Since $g(R_N) = p(R_L)$, therefore $i \in L$. Therefore, $R_L \succ R_M$, for all maximal peak-unanimous subprofiles R_M of R_N that are distinct from R_L . Therefore, g is the hierarchical 1-rule associated with the priority \succ , the desired conclusion. ■

Finally, we should point out that, although we defined strategy-proofness for 1-rules, the definition can easily be adapted for rules. It is easy to see that hierarchical rules are strategy-proof. Thus, our results in this section also imply that the only strategy-proof rules are the hierarchical rules.

We conclude with the following additional observations. All of the hierarchical rules satisfy hiding-proofness. However, the left-peaks rule and the right-peaks rule are the only anonymous hierarchical rules. Each majoritarian rule is hiding-proof and anonymous, but it is not a hierarchical rule, i.e. it is not strategy-proof. A necessary and sufficient condition for a hierarchical priority \succ to satisfy sovereignty is that the weak ordering \succeq of the agents that is

associated with it should have no maximal element.

10 Conclusion

In this paper we undertake the task of designing rules for the provision of multiple public facilities to populations of agents with single peaked preferences over the real line. To this end, we introduce two new axioms –*object-population monotonicity*– and *sovereignty*. We uncover and characterize a rich and interesting family of rules, the priority rules. Some rules in the family are appealing, for example the majoritarian rules. Unfortunately, no priority rule coincides with the median for each problem of providing exactly one public facility to an odd population of agents. On the positive side, the structure of the rules in the family is very simple, which makes it possible to characterize several subfamilies on the basis of anonymity, hiding-proofness and strategy-proofness. In particular, strategy-proof priority rules form the interesting subfamily of hierarchical rules.

We view our work as an initial step towards a comprehensive understanding of the implications of normative and strategic axioms in the problem of locating multiple facilities. An obvious question is whether it is possible to generalize Moulin’s (1980) generalized median voter schemes to this more complex setting. We only provide a partial answer to this question. The priority structure enables us to generalize only a particular type of generalized median voter scheme, the oligarchies. A generalization of the entire family of generalized median voter schemes to the problem of locating multiple facilities requires a flexibility which conflicts with the priority structure, and thus with our main axioms.

Our analysis is based on a particular way to extend preferences over single-locations to preferences over multiple locations, the max-extension, introduced by Miyagawa (1998, 2001). The work of Miyagawa seems to indicate that the max-extension is not very tractable, in the sense that it leads to characteriza-

tions of rules that are not particularly appealing, such as the left peaks rule and the right peaks rule, and other more complicated rules. As our results suggest, this may have more to do with the solidarity axioms used there than the max-extension itself. As Gordon (2007a) has shown, solidarity axioms have strong general implication in all public decisions models, of which the location of multiple public facilities is a special case. Nevertheless, both Miyagawa and us obtain highly discontinuous rules, for any natural topology. This seems to be a regular feature of the max-extension model. Thus, insisting on continuity in this model may be too strong a requirement. Majoritarian rules, for instance, are not continuous.

However, a natural extension of our work would be precisely to study what happens under other types of preference extensions, e.g. the lexicographic extension. Ehlers (2002, 2003) studies the same axioms as Miyagawa (1998, 2001) under the lexicographic extension and obtains strikingly different characterizations. So, it is quite possible that the axioms we used here would lead to other families of rules as well, if one were to replace the max-extension of our model with the lexicographic extension, or some other extension.

Finally, we have considered here the pure public goods case. Richer set-ups include problems with congestion and cost-sharing to finance the facilities, such as in the model by Jackson and Nicolò (2004), for the case $k = 1$. We leave this and the other questions open for future research.

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